

# Introduction: Microprogrammed Control(*Control Unit*)

## ■ Control Memory

### ◆ Control Unit

- Initiate sequences of microoperations
  - » Control signal (*that specify microoperations*) in a bus-organized system
    - groups of bits that select the paths in multiplexers, decoders, and arithmetic logic units
- Two major types of Control Unit
  - » Hardwired Control : *in Chap. 5*
    - The control logic is implemented with gates, F/Fs, decoders, and other digital circuits
    - + Fast operation, - Wiring change(if the design has to be modified)
  - » Microprogrammed Control : *in this Chapter*
    - The control information is stored in a control memory, and the control memory is programmed to initiate the required sequence of microoperations
    - + Any required change can be done by updating the microprogram in control memory, - Slow operation

### ◆ Control Word

- The control variables at any given time can be represented by a string of 1's and 0's.

### ◆ Microprogrammed Control Unit

- A control unit whose binary control variables are stored in memory (*control memory*).

## ◆ Microinstruction : *Control Word in Control Memory*

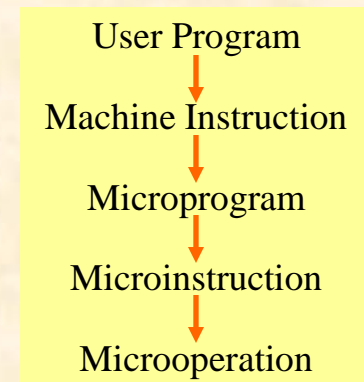
- The microinstruction specifies one or more microoperations

## ◆ Microprogram

- A sequence of microinstruction
  - » Dynamic microprogramming : *Control Memory = RAM*
    - RAM can be used for writing (*to change a writable control memory*)
    - Microprogram is loaded initially from an auxiliary memory such as a magnetic disk
  - » Static microprogramming : *Control Memory = ROM*
    - Control words in ROM are made permanent during the hardware production.

## ◆ Microprogrammed control Organization : *Fig. 7-1*

- 1) Control Memory
  - » A memory is part of a control unit : *Microprogram이 저장되어 있음*
  - » Computer Memory (*employs a microprogrammed control unit*)
    - Main Memory : for storing user program (*Machine instruction/data*)
    - Control Memory : for storing microprogram (*Microinstruction*)
- 2) Control Address Register
  - » Specify the address of the microinstruction
- 3) Sequencer (= *Next Address Generator*)
  - » Determine the address sequence that is read from control memory
  - » Next address of the next microinstruction can be specified several way depending on the sequencer input : *p. 217, [1, 2, 3, and 4]*



- 4) Control Data Register (= *Pipeline Register*)
  - » Hold the microinstruction read from control memory
  - » Allows the execution of the microoperations specified by the control word *simultaneously* with the generation of the next microinstruction

### ◆ RISC Architecture Concept

- RISC(Reduced Instruction Set Computer) system use hardwired control rather than microprogrammed control : **Sec. 8-8**

## ■ 2 Address Sequencing

### ◆ Address Sequencing = Sequencer : Next Address Generator

- Selection of address for control memory

### ◆ Routine → *Subroutine : program used by other ROUTINES*

- Microinstruction are stored in control memory *in groups*

### ◆ Mapping

- Instruction Code → Address in control memory(*where routine is located*)

### ◆ Address Sequencing Capabilities : *control memory address*

- 1) Incrementing of the control address register
- 2) Unconditional branch or conditional branch, depending on status bit conditions
- 3) Mapping process ( *bits of the instruction address for control memory* )
- 4) A facility for subroutine return

◆ Selection of address for control memory : *Fig. 7-2*

- Multiplexer

- ① CAR Increment

- ② JMP/CALL

- ③ Mapping

- ④ Subroutine Return

- CAR : Control Address Register

- » CAR receive the address from 4 different paths

- 1) Incrementer

- 2) Branch address from control memory

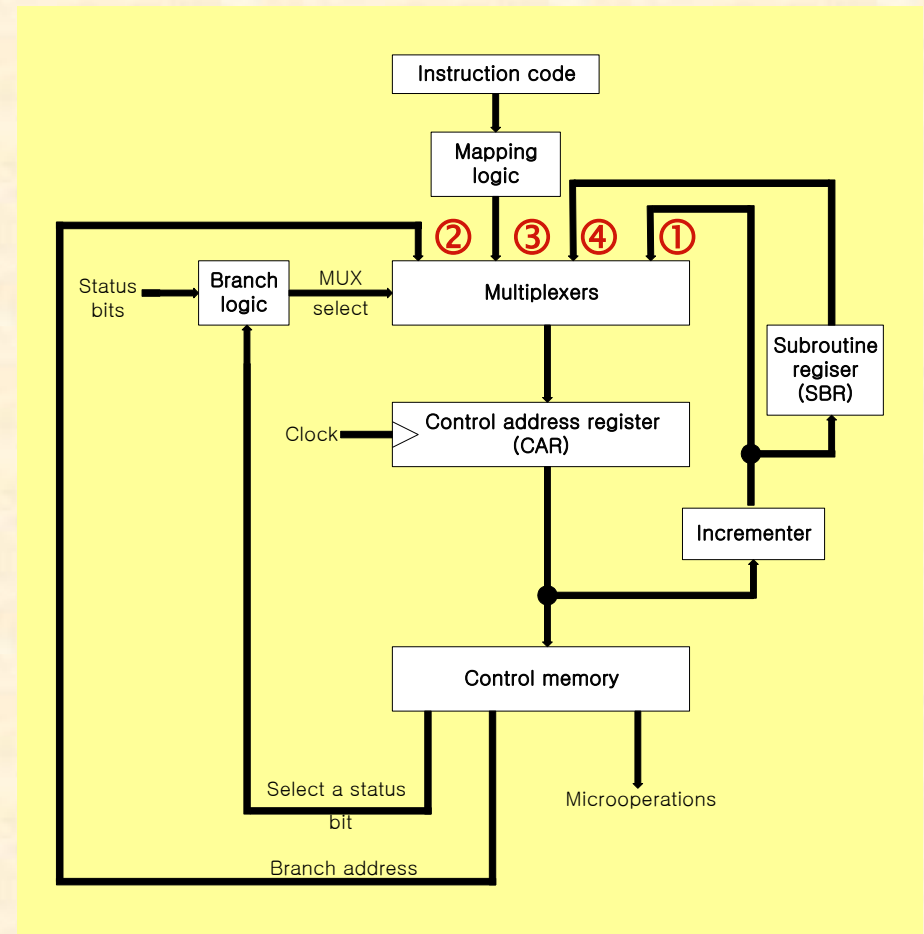
- 3) Mapping Logic

- 4) SBR : Subroutine Register

- SBR : Subroutine Register

- » Return Address can not be stored in ROM

- » Return Address for a subroutine is stored in SBR



## ◆ Conditional Branching

- Status Bits
  - » Control the conditional branch decisions generated in the **Branch Logic**
- Branch Logic
  - » Test the specified condition and Branch to the indicated address if the condition is met ; otherwise, the control address register is just incremented.
- Status Bit Test 와 Branch Logic의 실제 회로 : **Fig. 7-8**
  - » 4 X 1 Mux 와 Input Logic(**Tab. 7-4**)으로 구성

## ◆ Mapping of Instruction : **Fig. 7-3**

- Computer Instruction

Opcode				Address	
1	0	1	1		

### Mapping bits

0 x x x x 0 0

Microinstruction Address

0	1	0	1	1	0	0
---	---	---	---	---	---	---

- 4 bit Opcode = specify up to 16 distinct instruction
- Mapping Process : **Converts the 4-bit Opcode to a 7-bit control memory address**
  - » 1) Place a “0” in the most significant bit of the address
  - » 2) Transfer 4-bit Operation code bits
  - » 3) Clear the two least significant bits of the CAR ( 즉, 4 개의 Microinstruction 수용 가능 )
- Mapping Function : Implemented by **Mapping ROM** or **PLD**
- Control Memory Size : 128 words (=  $2^7$ )



## ◆ Subroutine

- Subroutines are programs that are used by other routines
  - » Subroutine can be called from any point within the main body of the microprogram
- Microinstructions can be saved by subroutines that use common section of microcode
  - » 예제) Memory Reference 명령에서 Operand의 Effective Address를 구하는 Subroutine
    - p. 228, Tab. 7-2에서 INDRCT (여기에서 FETCH와 INDRCT는 Subroutine)
    - Subroutine은 ORG 64, 즉 1000000 - 1111111에 위치(Routine은 0000000 - 0111111)
- Subroutine must have a provision for
  - » storing the return address during a subroutine call
  - » restoring the address during a subroutine return
    - Last-In First Out(LIFO) Register Stack : Sec. 8-7

## ■ 7-3 Microprogram Example

### ◆ Computer Configuration : Fig. 7-4

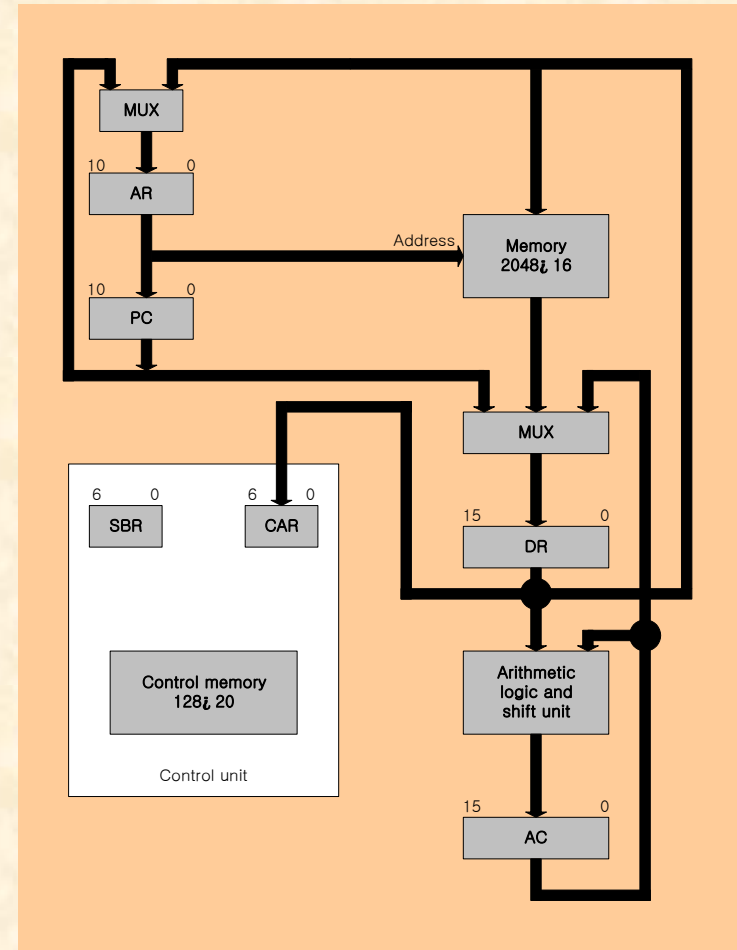
- 2 Memory : Main memory(instruction/data), Control memory(microprogram)
  - » Data written to memory come from DR, and Data read from memory can go only to DR
- 4 CPU Register and ALU : DR, AR, PC, AC, ALU
  - » DR can receive information from AC, PC, or Memory (selected by MUX)
  - » AR can receive information from PC or DR (selected by MUX)
  - » PC can receive information only from AR
  - » ALU performs microoperation with data from AC and DR (결과는 AC에 저장)
- 2 Control Unit Register : SBR, CAR

## ◆ Instruction Format

- Instruction Format : **Fig. 7-5(a)**
  - » I : 1 bit for indirect addressing
  - » Opcode : 4 bit operation code
  - » Address : 11 bit address for system memory
- Computer Instruction : **Fig. 7-5(b)**
  - » 16 명령어가 가능하며 4 개만 표시

## ◆ Microinstruction Format : **Fig. 7-6**

- 3 bit Microoperation Fields : F1, F2, F3
  - » 21 bit of Microoperation : **Tab. 7-1**
  - » 동시에 3 개의 microoperation 실행 가능
    - 3 개 이하일 경우, 000(no operation)으로 채움
  - » two or more conflicting microoperations can not be specified simultaneously
    - **예제)** 010 001 000  
Clear AC to 0 and subtract DR from AC at the same time
  - » Symbol **DRTAC**(F1 = 100)
    - stand for a transfer from DR to AC (*T = to*)



- 2 bit Condition Fields : CD
  - » 00 : Unconditional branch, **U** = 항상 1
  - » 01 : Indirect address bit, **I** = DR(15)
  - » 10 : Sign bit of AC, **S** = AC(15)
  - » 11 : Zero value in AC, **Z** = AC = 0
- 2 bit Branch Fields : BR
  - » 00 : **JMP**
    - Condition = 0 : 1  $CAR \leftarrow CAR + 1$
    - Condition = 1 : 2  $CAR \leftarrow AD$
  - » 01 : **CALL**
    - Condition = 0 : 1  $CAR \leftarrow CAR + 1$
    - Condition = 1 : 2  $CAR \leftarrow AD$ , 3  $SBR \leftarrow CAR + 1$
  - » 10 : **RET** 3  $CAR \leftarrow SBR$
  - » 11 : **MAP** 4  $CAR(2-5) \leftarrow DR(11-14), CAR(0, 1, 6) \leftarrow 0$
- 7 bit Address Fields : AD
  - » 128 word : 128 X 20 bit

Save Return Address

Restore Return Address

## ◆ Symbolic Microinstruction

- ① Label Field : Terminated with a colon ( : )
- ② Microoperation Field : one, two, or three symbols, separated by commas
- ③ CD Field : U, I, S, or Z
- ④ BR Field : JMP, CALL, RET, or MAP

① Label	② Microoperat	③ CD	④ BR	⑤ AD
FETCH:	ORG 64			
	PCTAR	U	JMP	NEXT
	READ, INCPC	U	JMP	NEXT
	DRTAR	U	MAP	0



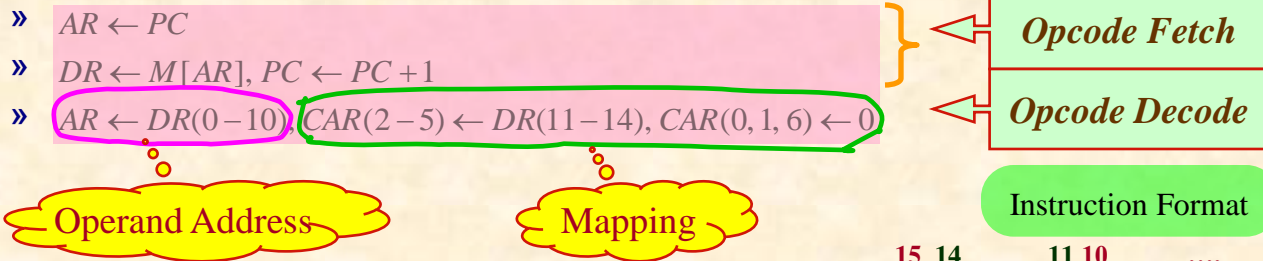
## ⑤ AD Field

- Symbolic Address : Label ( = Address )
- Symbol "NEXT" : next address
- Symbol "RET" or "MAP" : AD field = 0000000

- ORG : Pseudoinstruction(*define the origin, or first address of routine*)

## ◆ Fetch (Sub)Routine

- Memory Map(128 words) : **Tab. 7-2, Tab. 7-3**
  - » Address 0 to 63 : Routines for the 16 instruction(현재는 4 instruction)
  - » Address 64 to 127 : Any other purpose(현재는 Subroutines : FETCH, INDRCT)
- Microinstruction for FETCH Subroutine



- Fetch Subroutine : address 64

Label	Microoperat	CD	BR	AD
	ORG 64			
FETCH:	PCTAR	U	JMP	NEXT
	READ, INCPC	U	JMP	NEXT
	DRTAR	U	MAP	0

## ◆ Symbolic Microprogram : *Tab. 7-2*

### ● The execution of MAP microinstruction in FETCH subroutine

#### » Branch to address $0xxxx00$ ( $xxxx = 4 \text{ bit Opcode}$ )

- ADD :  $0 \text{ } 0000 \text{ } 00 = 0$
- BRANCH :  $0 \text{ } 0001 \text{ } 00 = 4$
- STORE :  $0 \text{ } 0010 \text{ } 00 = 8$
- EXCHANGE :  $0 \text{ } 0011 \text{ } 00 = 12, (16, 20, \dots, 60)$

### ● Indirect Address : $I = 1$

#### » Indirect Addressing : $AR \leftarrow M[AR]$

- AR이 지시하는 메모리 내용을 DR에 읽은 후, 다시 AR에 써 넣는다

#### » INDRCT subroutine

Label	Microoperat	CD	BR	AD	
INDRCT:	READ	U	JMP	NEXT	} → $DR \leftarrow M[AR]$ $AR \leftarrow DR$
	DRTAR	U	RET	0	

### ● Execution of Instruction

#### » ADD instruction 실행 절차

- 1) ADD 명령이 실행되면 FETCH subroutine에서 Opcode를 fetch한 후, MAP이 실행되면 MAP Process에 의해  $CAR = 0 \text{ } 0000 \text{ } 00$ 으로 branch 한다( 여기서  $Opcode = 0000$ , *Fig. 7-5(b)* )
- 2) ADD 명령의 Address 0 에서는 CD 비트를 검사하여 Indirect = 1이면 INDRCT subroutine에서 Effective Address를 AR에 써넣고 Return 한다.
- 3) ADD 명령의 Address 1 에서는 AR이 지시하는 Memory의 내용을 읽어서 DR에 써넣는다.
- 4) ADD 명령의 Address 2 에서는  $AC + DR$ 을 AC에 저장한 후, FETCH subroutine으로 Branch하면 1)에서와 같은 방법으로 PC가 지시하는 명령어를 Fetch 하여 MAP 수행 결과에 따라 해당 Routine Address로 Branch 한다.

» BRANCH instruction 실행 절차

- 1) BRANCH 명령의 Address 4 에서는 CD Bit를 검사하여 Sign(S) = 1 이면 Address 6 번으로 가서 Indirect를 검사하고 ARTPC에 의해 해당 Address로 Branch 한 후, FETCH에 의해 PC 가 지시하는 다음 명령을 수행한다.
- 2) BRANCH 명령의 Address 4에서 Sign = 0 이면 Branch 하지 않고 FETCH에 의해 PC가 지시하는 다음 명령을 수행한다.

» STORE instruction 실행 절차

» EXCHANGE instruction 실행 절차

◆ Binary Microprogram : *Tab. 7-3*

- Symbolic microprogram(*Tab. 7-2*) must be translated to **binary** either by means of an assembler program or by the user
- Control Memory
  - » Most microprogrammed systems use a ROM for the control memory
    - Cheaper and faster than a RAM
    - Prevent the occasional user from changing the architecture of the system

■ 7-4 Design of Control Unit

◆ Decoding of Microinstruction Fields : *Fig. 7-7*

- F1, F2, and F3 of Microinstruction are decoded with a 3 x 8 decoder
- Output of decoder must be connected to the proper circuit to initiate the corresponding microoperation (*as specified in Tab. 7-1*)

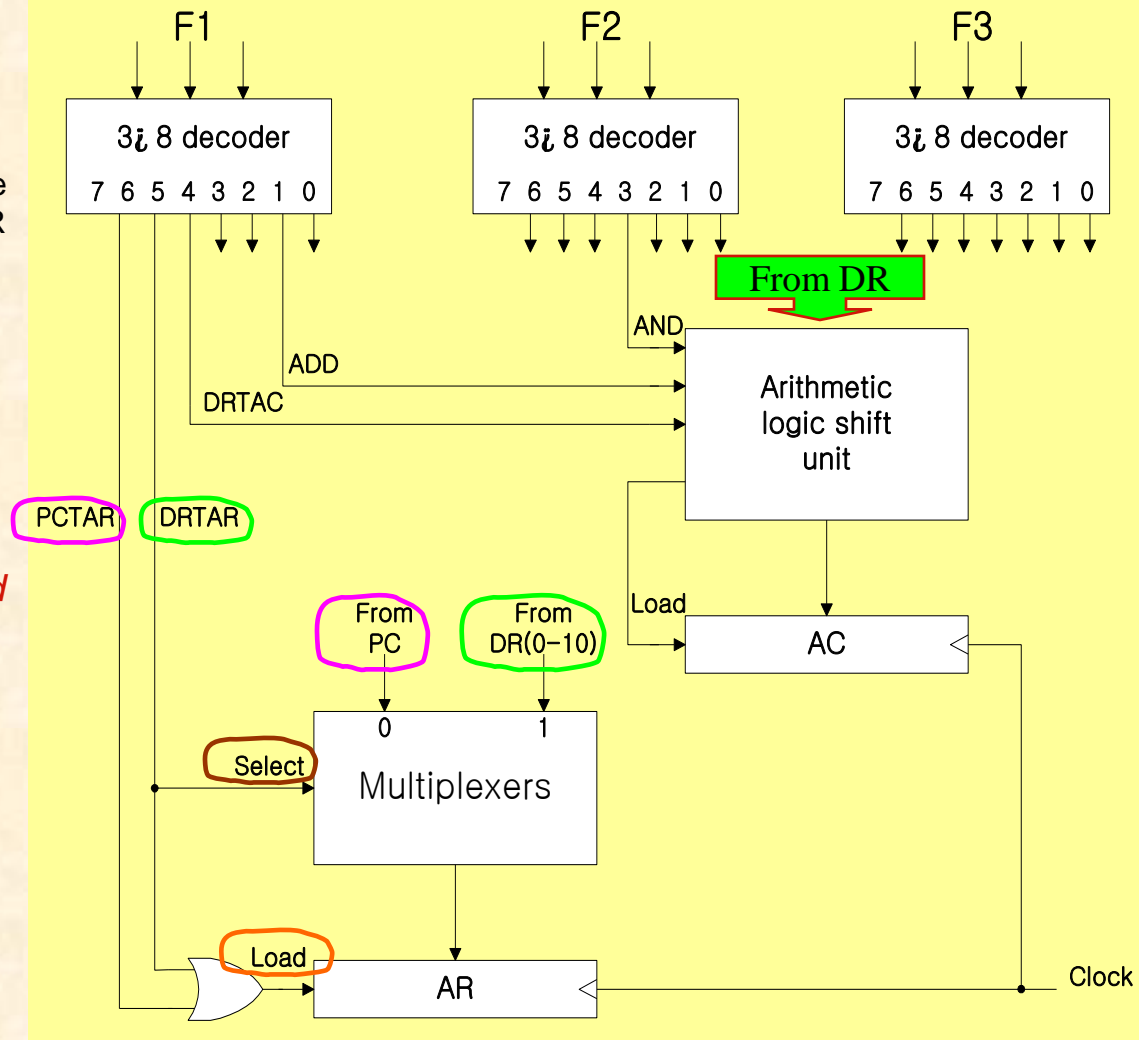
» 예제) F1 = 101 (5) : **DRTAR**

F1 = 110 (6) : **PCTAR**

- Output 5 and 6 of decoder F1 are connected to the load input of AR (two input of OR gate)
- Multiplexer select the data from DR when output 5 is active
- Multiplexer select the data from AC when output 5 is inactive

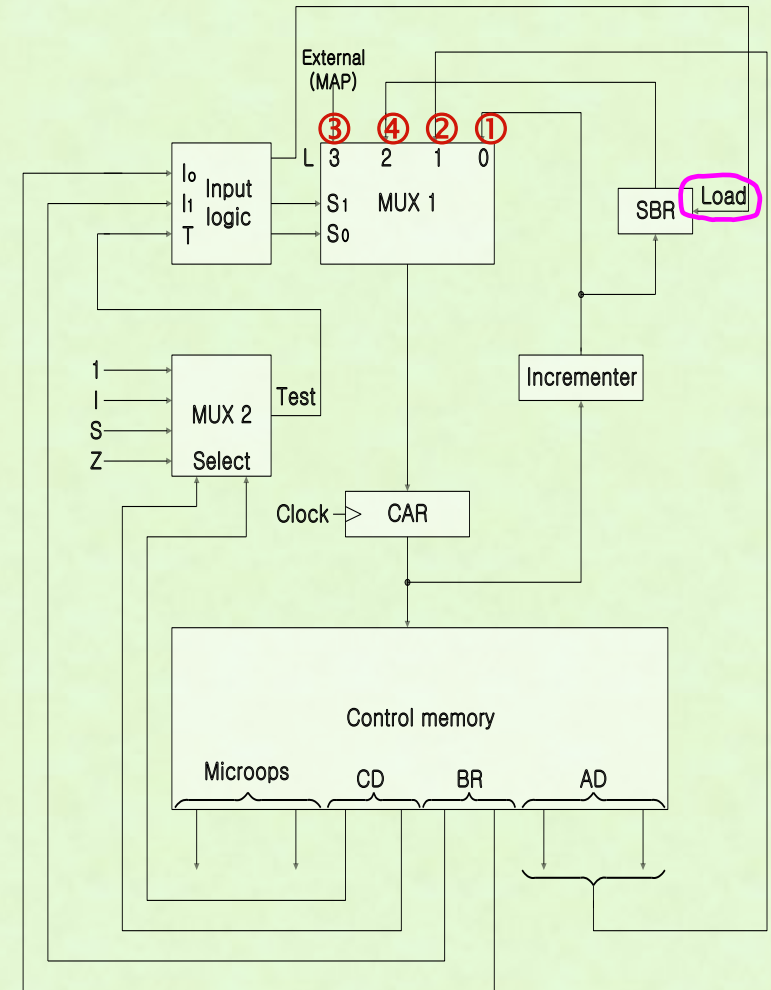
## ● Arithmetic Logic Shift Unit

- » Control signal of ALU in *hardwired control*: p. 164, Fig. 5-19, 20
- » Control signal will be now come from the *output of the decoders* associated with the AND, ADD, and DRTAC.



## ◆ Microprogram Sequencer : *Fig. 7-8*

- Microprogram Sequencer select the next address for control memory
- MUX 1
  - » Select an address source and route to CAR
    - ① CAR + 1
    - ② JMP/CALL
    - ③ Mapping
    - ④ Subroutine Return
  - » JMP 와 CALL의 차이점
    - JMP : AD가 MUX 1의 2번을 통해 CAR로 전송
    - CALL : AD가 MUX 1의 2번을 통해 CAR로 전송되고, 동시에 CAR + 1(Return Address) 이 LOAD 신호에 의해 SBR에 저장된다.
- MUX 2
  - » Test a status bit and the result of the test is applied to an input logic circuit
  - » One of 4 Status bit is selected by Condition bit (CD)
- Design of Input Logic Circuit
  - » Select one of the source address( $S_0$ ,  $S_1$ ) for CAR
  - » Enable the load input(L) in SBR





- Input Logic Truth Table : **Tab. 7-4**

- » Input :

- $I_0, I_1$  from Branch bit (**BR**)
- T from MUX 2 (**T**)

- » Output :

- MUX 1 Select signal (**S<sub>0</sub>, S<sub>1</sub>**)  
 $S1 = I_1 I_0' + I_1 I_0 = I_1 (I_0' + I_0) = I_1$   
 $S0 = I_1' I_0' T + I_1' I_0 T + I_1 I_0$   
 $= I_1' T (I_0' + I_0) + I_1 I_0$   
 $= I_1' T + I_1 I_0$
- SBR Load signal (**L**)  
 $L = I_1' I_0 T$

BR Field		Input			MUX 1		Load SBR
		I1	I0	T	S1	S0	L
0	0	0	0	0	0	0	0
0	0	0	0	1	0	1	0
0	1	0	1	0	0	0	0
0	1	0	1	1	0	1	1
1	0	1	0	x	1	0	0
1	1	1	1	x	1	1	0

- ① CAR + 1
- ② JMP
- ① CAR + 1
- ② CALL
- ③ MAP
- ④ RET

**CALL**

# Assignment

- What do you mean by control unit.
- Explain micro programmed control unit.
- What do you mean by micro sequencing.